

Frames and a vector-valued ambiguity function

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- 1 Problem and goal
- 2 Frames
- 3 Multiplication problem and A_p^1
- 4 $A_p^d : \mathbb{Z}_N \times \mathbb{Z}_N \longrightarrow \mathbb{C}^d, u : \mathbb{Z}_N \longrightarrow \mathbb{C}^d$
- 5 $A_p^d(u)$ for DFT frames
- 6 Figure
- 7 Epilogue

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- 2 Frames
- 3 Multiplication problem and A_p^1
- 4 $A_p^d : \mathbb{Z}_N \times \mathbb{Z}_N \longrightarrow \mathbb{C}^d, u : \mathbb{Z}_N \longrightarrow \mathbb{C}^d$
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- 6 Figure
- 7 Epilogue

- Originally, our problem was to construct libraries of phase-coded waveforms v parameterized by design variables, for communications and radar.
- A goal was to achieve diverse ambiguity function behavior of v by defining new classes of quadratic phase and number theoretic perfect autocorrelation codes u with which to define v .
- A realistic more general problem was to construct vector-valued waveforms v in terms of vector-valued perfect autocorrelation codes u . Such codes are relevant in light of vector sensor and MIMO capabilities and modeling.
- Example: Discrete time data vector $u(k)$ for a d -element array,

$$k \mapsto u(k) = (u_0(k), \dots, u_{d-1}(k)) \in \mathbb{C}^d.$$

We can have $\mathbb{R}^N \rightarrow GL(d, \mathbb{C})$, or even more general.

General problem and STFT theme

- Establish the theory of vector-valued ambiguity functions to estimate v in terms of ambiguity data.
- First, establish this estimation theory by defining the discrete periodic vector-valued ambiguity function in a natural way.
- Mathematically, this natural way is to formulate the discrete periodic vector-valued ambiguity function in terms of the Short Time Fourier Transform (STFT).

STFT and ambiguity function

Short time Fourier transform – STFT

- The narrow band cross-correlation ambiguity function of v, w defined on \mathbb{R} is

$$A(v, w)(t, \gamma) = \int_{\mathbb{R}} v(s+t) \overline{w(s)} e^{-2\pi i s \gamma} ds.$$

- $A(v, w)$ is the STFT of v with window w .
- The *narrow band radar ambiguity function* $A(v)$ of v on \mathbb{R} is

$$\begin{aligned} A(v)(t, \gamma) &= \int_{\mathbb{R}} v(s+t) \overline{v(s)} e^{-2\pi i s \gamma} ds \\ &= e^{\pi i t \gamma} \int_{\mathbb{R}} v\left(s + \frac{t}{2}\right) \overline{v\left(s - \frac{t}{2}\right)} e^{-2\pi i s \gamma} ds, \text{ for } (t, \gamma) \in \mathbb{R}^2. \end{aligned}$$

Goal

- Let v be a phase coded waveform with N lags defined by the code u .
- Let u be N -periodic, and so $u : \mathbb{Z}_N \rightarrow \mathbb{C}$, where \mathbb{Z}_N is the additive group of integers modulo N .
- The *discrete periodic ambiguity function* $A_p(u) : \mathbb{Z}_N \times \mathbb{Z}_N \rightarrow \mathbb{C}$ is

$$A_p(u)(m, n) = \frac{1}{N} \sum_{k=0}^{N-1} u(m+k) \overline{u(k)} e^{-2\pi i kn/N}.$$

Goal

Given a vector valued N -periodic code $u : \mathbb{Z}_N \rightarrow \mathbb{C}^d$, construct the following in a meaningful, computable way:

- Generalized \mathbb{C} -valued periodic ambiguity function $A_p^1(u) : \mathbb{Z}_N \times \mathbb{Z}_N \rightarrow \mathbb{C}$
- \mathbb{C}^d -valued periodic ambiguity function $A_p^d(u) : \mathbb{Z}_N \times \mathbb{Z}_N \rightarrow \mathbb{C}^d$

The STFT is the *guide* and the *theory of frames* is the technology to obtain the goal.

- 1 Problem and goal
- 2 **Frames**
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- 5 $A_p^d(u)$ for DFT frames
- 6 Figure
- 7 Epilogue

- A sequence $F = \{E_j\}_{j=1}^N \subseteq \mathbb{C}^d$ is a *frame* for \mathbb{C}^d if

$$\exists A, B > 0 \quad \text{such that} \quad \forall u \in \mathbb{C}^d, \quad A\|u\|^2 \leq \sum_{j=1}^N |\langle u, E_j \rangle|^2 \leq B\|u\|^2.$$

- F is a *tight frame* if $A = B$; and F is a *finite unit-norm tight frame* (FUNTF) if $A = B$ and each $\|E_j\| = 1$.
- Theorem: If $\{E_j\}_{j=0}^{N-1}$ is a FUNTF for \mathbb{C}^d , then

$$\forall u \in \mathbb{C}^d, \quad u = \frac{d}{N} \sum_{j=0}^{N-1} \langle u, E_j \rangle E_j.$$

- Frames give redundant signal representation to compensate for hardware errors, to ensure numerical stability, and to minimize the effects of noise.

Recent applications of FUNTFs

- Robust transmission of data over erasure channels such as the internet [Casazza, Goyal, Kelner, Kovačević]
- Multiple antenna code design for wireless communications [Hochwald, Marzetta, T. Richardson, Sweldens, Urbanke]
- Multiple description coding [Goyal, Heath, Kovačević, Strohmer, Vetterli]
- Quantum detection [Bölcskei, Eldar, Forney, Oppenheim, Kebo, B]
- Grassmannian "min-max" waveforms [Calderbank, Conway, Sloane, et al., Kolesar, B]

Let $N \geq d$ and form an $N \times d$ matrix using any d columns of the $N \times N$ DFT matrix $(e^{2\pi ijk/N})_{j,k=0}^{N-1}$. The rows of this $N \times d$ matrix, up to multiplication by $\frac{1}{\sqrt{d}}$, form a FUNTF for \mathbb{C}^d .

$$N = 8, d = 5 \quad \frac{1}{\sqrt{5}} \begin{bmatrix} * & * & \cdot & \cdot & * & * & * & \cdot \\ * & * & \cdot & \cdot & * & * & * & \cdot \\ * & * & \cdot & \cdot & * & * & * & \cdot \\ * & * & \cdot & \cdot & * & * & * & \cdot \\ * & * & \cdot & \cdot & * & * & * & \cdot \\ * & * & \cdot & \cdot & * & * & * & \cdot \\ * & * & \cdot & \cdot & * & * & * & \cdot \\ * & * & \cdot & \cdot & * & * & * & \cdot \end{bmatrix}$$

$$x_m = \frac{1}{5} (e^{2\pi i \frac{m}{8}}, e^{2\pi i m \frac{2}{8}}, e^{2\pi i m \frac{5}{8}}, e^{2\pi i m \frac{6}{8}}, e^{2\pi i m \frac{7}{8}})$$

$$m = 1, \dots, 8.$$

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- 5 $A_p^d(u)$ for DFT frames
- 6 Figure
- 7 Epilogue

Multiplication problem

- Given $u : \mathbb{Z}_N \rightarrow \mathbb{C}^d$.
- If $d = 1$ and $e_n = e^{2\pi i n/N}$, then

$$A_p(u)(m, n) = \frac{1}{N} \sum_{k=0}^{N-1} \langle u(m+k), u(k) e_{nk} \rangle.$$

Multiplication problem

To characterize sequences $\{E_k\} \subseteq \mathbb{C}^d$ and multiplications $*$ so that

$$A_p^1(u)(m, n) = \frac{1}{N} \sum_{k=0}^{N-1} \langle u(m+k), u(k) * E_{nk} \rangle \in \mathbb{C}$$

is a meaningful and well-defined *ambiguity function*. This formula is clearly motivated by the STFT.

Ambiguity function assumptions

There is a natural way to address the multiplication problem motivated by the fact that $e_m e_n = e_{m+n}$. To this end, we shall make the *ambiguity function assumptions*:

- $\exists \{E_k\}_{k=0}^{N-1} \subseteq \mathbb{C}^d$ and a multiplication $*$ such that $E_m * E_n = E_{m+n}$ for $m, n \in \mathbb{Z}_N$;
- $\{E_k\}_{k=0}^{N-1} \subseteq \mathbb{C}^d$ is a tight frame for \mathbb{C}^d ;
- $*$: $\mathbb{C}^d \times \mathbb{C}^d \rightarrow \mathbb{C}^d$ is bilinear, in particular,

$$\left(\sum_{j=0}^{N-1} c_j E_j \right) * \left(\sum_{k=0}^{N-1} d_k E_k \right) = \sum_{j=0}^{N-1} \sum_{k=0}^{N-1} c_j d_k E_j * E_k.$$

- Let $\{E_j\}_j^{N-1} \subseteq \mathbb{C}^d$ satisfy the three ambiguity function assumptions.
- Given $u, v : \mathbb{Z}_N \longrightarrow \mathbb{C}^d$ and $m, n \in \mathbb{Z}_N$.
- Then, one calculates

$$u(m) * v(n) = \frac{d^2}{N^2} \sum_{j=0}^{N-1} \sum_{s=0}^{N-1} \langle u(m), E_j \rangle \langle v(n), E_s \rangle E_{j+s}.$$

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- 2 Frames
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- 4 $A_p^d : \mathbb{Z}_N \times \mathbb{Z}_N \longrightarrow \mathbb{C}^d, u : \mathbb{Z}_N \longrightarrow \mathbb{C}^d$**
- 5 $A_p^d(u)$ for DFT frames
- 6 Figure
- 7 Epilogue

$A_p^1(u)$ for DFT frames

- Let $\{E_j\}_{j=0}^{N-1} \subseteq \mathbb{C}^d$ satisfy the three ambiguity function assumptions.
- Further, assume that $\{E_j\}_{j=0}^{N-1}$ is a DFT frame, and let r designate a fixed column.
- Without loss of generality, choose the first d columns of the $N \times N$ DFT matrix.
- Then, one calculates

$$\begin{aligned} E_m * E_n(r) &= \frac{d^2}{N^2} \sum_{j=0}^{N-1} \sum_{s=0}^{N-1} \langle E_m, E_j \rangle \langle E_n, E_s \rangle E_{j+s}(r). \\ &= \frac{e^{(m+n)r}}{\sqrt{d}} = E_{m+n}(r). \end{aligned}$$

$A_p^1(u)$ for DFT frames

- Thus, for DFT frames, $*$ is componentwise multiplication in \mathbb{C}^d with a factor of \sqrt{d} .
- In this case $A_p^1(u)$ is well-defined for $u : \mathbb{Z}_N \longrightarrow \mathbb{C}^d$ by

$$\begin{aligned} A_p^1(u)(m, n) &= \frac{1}{N} \sum_{k=0}^{N-1} \langle u(m+k), u(k) * E_{nk} \rangle \\ &= \frac{d}{N^2} \sum_{k=0}^{N-1} \sum_{j=0}^{N-1} \langle E_j, u(k) \rangle \langle u(m+k), E_{j+nk} \rangle. \end{aligned}$$

- In the previous DFT example, $*$ is intrinsically related to the “addition” defined on the indices of the frame elements, viz.,
 $E_m * E_n = E_{m+n}$.
- Alternatively, we could have $E_m * E_n = E_{m \bullet n}$ for some function $\bullet : \mathbb{Z}_N \times \mathbb{Z}_N \longrightarrow \mathbb{Z}_N$, and, thereby, we could use frames which are not FUNTFs.
- Given a bilinear multiplication $* : \mathbb{C}^d \times \mathbb{C}^d \longrightarrow \mathbb{C}^d$, we can find a frame $\{E_j\}_j$ and an index operation \bullet with the $E_m * E_n = E_{m \bullet n}$ property.
- If \bullet is the multiplication for a group, possibly non-abelian and/or infinite, we may reverse the process and find a FUNTF and bilinear multiplication $*$ with the $E_m * E_n = E_{m \bullet n}$ property.

$A_p^1(u)$ for cross product frames

- Take $* : \mathbb{C}^3 \times \mathbb{C}^3 \longrightarrow \mathbb{C}^3$ to be the cross product on \mathbb{C}^3 and let $\{i, j, k\}$ be the standard basis.
- $i * j = k, j * i = -k, k * i = j, i * k = -j, j * k = i, k * j = -i,$
 $i * i = j * j = k * k = 0.$ $\{0, i, j, k, -i, -j, -k, \}$ is a tight frame for \mathbb{C}^3 with frame constant 2. Let
 $E_0 = 0, E_1 = i, E_2 = j, E_3 = k, E_4 = -i, E_5 = -j, E_6 = -k.$
- The index operation corresponding to the frame multiplication is the non-abelian operation $\bullet : \mathbb{Z}_7 \times \mathbb{Z}_7 \longrightarrow \mathbb{Z}_7,$ where
 $1 \bullet 2 = 3, 2 \bullet 1 = 6, 3 \bullet 1 = 2, 1 \bullet 3 = 5, 2 \bullet 3 = 1, 3 \bullet 2 = 4, 1 \bullet 1 =$
 $2 \bullet 2 = 3 \bullet 3 = 0, n \bullet 0 = 0 \bullet n = 0, 1 \bullet 4 = 0, 1 \bullet 5 = 6, 1 \bullet 6 = 2, 4 \bullet 1 =$
 $0, 5 \bullet 1 = 3, 6 \bullet 1 = 5, 2 \bullet 4 = 3, 2 \bullet 5 = 0,$ etc.
- The three ambiguity function assumptions are valid and so we can write the cross product as

$$u \times v = u * v = \frac{1}{2^2} \sum_{s=1}^6 \sum_{t=1}^6 \langle u, E_s \rangle \langle v, E_t \rangle E_{s \bullet t}.$$

- Consequently, $A_p^1(u)$ can be well-defined.

Vector-valued ambiguity function $A_p^d(u)$

- Let $\{E_j\}_j^{N-1} \subseteq \mathbb{C}^d$ satisfy the three ambiguity function assumptions.
- Given $u : \mathbb{Z}_N \rightarrow \mathbb{C}^d$.
- The following definition is clearly *motivated* by the STFT.

Definition

$A_p^d(u) : \mathbb{Z}_N \times \mathbb{Z}_N \rightarrow \mathbb{C}^d$ is defined by

$$A_p^d(u)(m, n) = \frac{1}{N} \sum_{k=0}^{N-1} u(m+k) * \overline{u(k)} * \overline{E_{nk}}.$$

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- 2 Frames
- 3 Multiplication problem and A_p^1
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- 6 Figure
- 7 Epilogue

STFT formulation of $A_p(u)$

- The discrete periodic ambiguity function of $u : \mathbb{Z}_N \rightarrow \mathbb{C}$ can be written as

$$A_p(u)(m, n) = \frac{1}{N} \sum_{k=0}^{N-1} \langle \tau_m u(k), F^{-1}(\tau_n \hat{u})(k) \rangle,$$

where $\tau(m)u(k) = u(m+k)$ is translation by m and $F^{-1}(u)(k) = \check{u}(k)$ is Fourier inversion.

- As such we see that $A_p(u)$ has the form of a STFT.
- We shall develop a vector-valued DFT theory to *verify* (not just *motivate*) that $A_p^d(u)$ is an STFT in the case $\{E_k\}_{k=0}^{N-1}$ is a DFT frame for \mathbb{C}^d .

Definition

Given $u : \mathbb{Z}_N \rightarrow \mathbb{C}^d$, and let $\{E_k\}_{k=0}^{N-1}$ be a DFT frame for \mathbb{C}^d . The *vector-valued discrete Fourier transform* of u is

$$\forall n \in \mathbb{Z}_N, \quad F(u)(n) = \hat{u}(n) = \sum_{m=0}^{N-1} u(m) * E_{mn},$$

where $*$ is pointwise (coordinatewise) multiplication.

Vector-valued Fourier inversion theorem

- Inversion process for the vector-valued case is analogous to the 1-dimensional case.
- We must define a new multiplication in the frequency domain to avoid divisibility problems.
- Define the weighted multiplication $(*) : \mathbb{C}^d \times \mathbb{C}^d \rightarrow \mathbb{C}^d$ by $u(*)v = u * v * \omega$ where $\omega = (\omega_1, \dots, \omega_d)$ has the property that each $\omega_n = \frac{1}{\#\{m \in \mathbb{Z}_N : mn=0\}}$.
- For the following theorem assume $d \ll N$ or N prime.

Theorem - Vector-valued Fourier inversion

The vector valued Fourier transform F is an isomorphism from $\ell^2(\mathbb{Z}_N)$ to $\ell^2(\mathbb{Z}_N, \omega)$ with inverse

$$\forall m \in \mathbb{Z}_N, \quad F^{-1}(m) = u(m) = \frac{d}{N} \sum_{n=0}^{N-1} \hat{u}(n) * E_{-mn} * \omega.$$

N prime implies F is unitary.

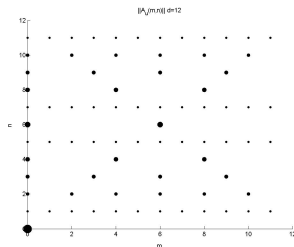
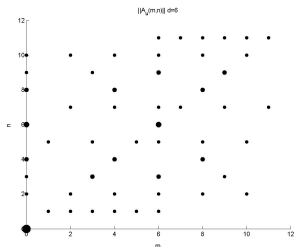
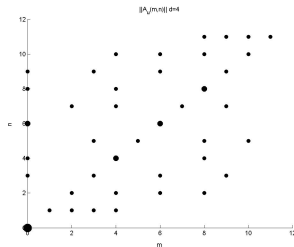
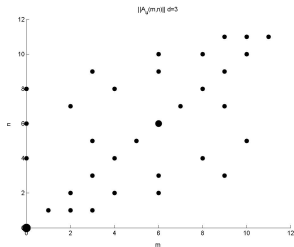
- Given $u, v : \mathbb{Z}_N \longrightarrow \mathbb{C}^d$, and let $\{E_k\}_{k=0}^{N-1}$ be a DFT frame for \mathbb{C}^d .
- $u * \bar{v}$ denotes pointwise (coordinatewise) multiplication with a factor of \sqrt{d} .
- We compute

$$A_p^d(u)(m, n) = \frac{1}{N} \sum_{k=0}^{N-1} (\tau_m u(k)) * \overline{F^{-1}(\tau_n \hat{u})(k)}.$$

- Thus, $A_p^d(u)$ is compatible with point of view of defining a vector-valued ambiguity function in the context of the STFT.

- 1 Problem and goal
- 2 Frames
- 3 Multiplication problem and A_p^1
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- 6 Figure**
- 7 Epilogue

Intensity $\|A_u^d(m, n)\|$ for $d = 3, 4, 6, 12$ where $u : \mathbb{Z}_N \rightarrow \mathbb{C}^d$ is a Wiener CAZAC



Outline

- 1 Problem and goal
- 2 Frames
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- 6 Figure
- 7 Epilogue

- If (G, \bullet) is a finite group with representation $\rho : G \rightarrow GL(\mathbb{C}^d)$, then there is a frame $\{E_n\}_{n \in G}$ and bilinear multiplication, $*$: $\mathbb{C}^d \times \mathbb{C}^d \rightarrow \mathbb{C}^d$, such that $E_m * E_n = E_{m \bullet n}$. Thus, we can develop $A_p^d(u)$ theory in this setting.
- Analyze ambiguity function behavior for (phase-coded) vector-valued waveforms $v : \mathbb{R} \rightarrow \mathbb{C}^d$, defined by $u : \mathbb{Z}_N \rightarrow \mathbb{C}^d$ as

$$v = \sum_{k=0}^{N-1} u(k) \mathbb{1}_{[kT, (k+1)T)},$$

in terms of $A_p^d(u)$. (See Figure)

That's all folks!